

Local-area network technologies

Shared-medium broadcasting

LAN topologies

802 MACs & Ethernet

LANs & TCP/IP

ARP & DHCP

Non-TCP/IP LANs

Overview

- Compared to WANs, local-area networks:
 - serve small areas (reduced geographic scope)
 - serve desktops and departmental servers
 - are generally connected by high-speed communications channels
- Varying philosophies, differing technologies
 - switched point-to-point
 - shared-medium packet broadcasting
- Data link and physical layers
- References:
 - Tanenbaum: Ch1.2.1, 4.1, 4.2.1–4.2.3, 4.3, 4.5.1, 4.5.2, 5.5
 - Kurose & Ross: Ch 5.1, 5.3, 5.5–5.7, 5.4

Switched point-to-point

- Evolution of telephone technology:
 - PBX – private branch exchange
 - digital PBX
 - digital data switch
- Based on a “switch” (hardware or software or hybrid) that creates point-to-point connections
- The distinction between telephony and data communications is blurred
 - telco manufacturers provide “data options” for digital PBXs
 - networking manufacturers creating voice capabilities for data networks

Shared-medium broadcasting

- In shared-medium broadcasting, all hosts transmit & receive in the same medium (channel)
 - broadcasting: every host receives every frame sent by any host
 - channel is *multiple access*
 - analogous to two-way radio with a single frequency
 - this environment requires careful coordination: handled by *Medium Access Control* (MAC) protocols (MAC layer)
 - *contention*: what happens if more than one host wants to transmit?
 - transmission *collisions* might occur: electrical interference that causes each transmissions to be garbled and unusable
 - how to deal with contention
 - by complete avoidance
 - by retransmission (after collision)

... shared-medium broadcasting, 2

- Basic techniques:
 - MAC layer manages *channel allocation*
 - two basic styles: static vs dynamic
 - for static channels, use TDM or FDM to divide the channel in advance
 - class of protocols called *partitioning protocols*
 - suitable for busy networks or networks that require predictable performance, but generally wasteful for non-continuous traffic
 - for dynamic allocation, hosts used entire channel “on demand” whenever they want
 - protocols classified as *random access* or *taking-turn*
 - taking-turn protocols avoid collisions
 - random-access protocols allow hosts to transmit at random and collide
 - collisions are handled by retransmission

Broadcast: contention

- Various methods of resolving contention
- ALOHA:
 - transmit at random, then listen to determine collision
 - wait and retransmit if necessary
- Carrier sense:
 - listen to channel
 - if channel is idle, start transmitting
 - if channel is busy, wait a while and retry
- Carrier sense with collision detection:
 - listen, transmit if channel is available
 - listen during transmission
 - stop transmitting immediately if collision
 - collisions can still occur because of the latency of electrical signals
 - wait and retransmit

Topology and media

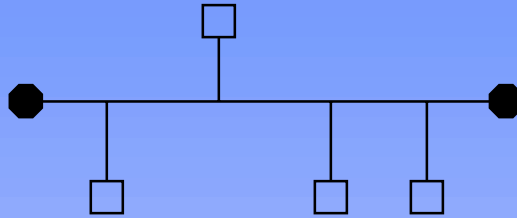
- Topology: the logical (as opposed to physical) arrangement of the hosts
 - two types:
 - bus (also called tree)
 - ring
 - star is sometimes considered a topology
- Choices for medium:
 - twisted-pair
 - baseband coaxial
 - broadband coaxial
 - fiber-optic
 - wireless

... topology and media, 2

- In theory, any medium for any topology
 - in practice, only a few combinations are used
 - twisted pair dominates desktop and most installations
 - fiber used for MAN, WAN and backbones (fiber to the desktop still on the horizon)
 - baseband coaxial is legacy only
 - broadband coaxial is used as part of cableco HFC, and in special-purpose industrial applications
- Wireless is an emerging technology
 - wireless LANs for building interiors to eliminate wiring costs & mobile computers
 - wireless for LANs is separate from wireless for hand-held PDAs & similar devices in the cellular world
 - wireless is also used for WAN point-to-point (microwave & related) which is also not related to wireless LANs

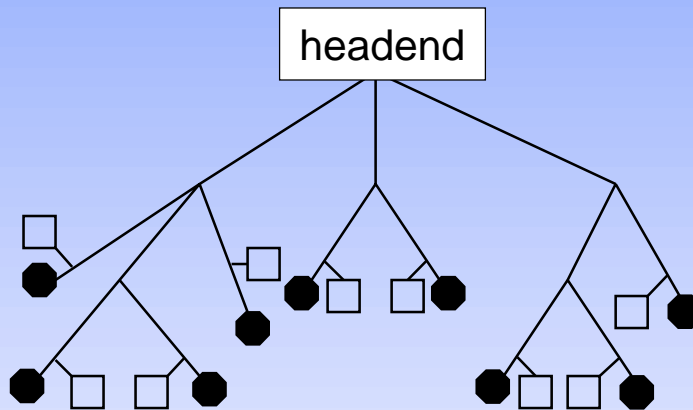
Bus & tree

- Bus:



- data moves in both directions

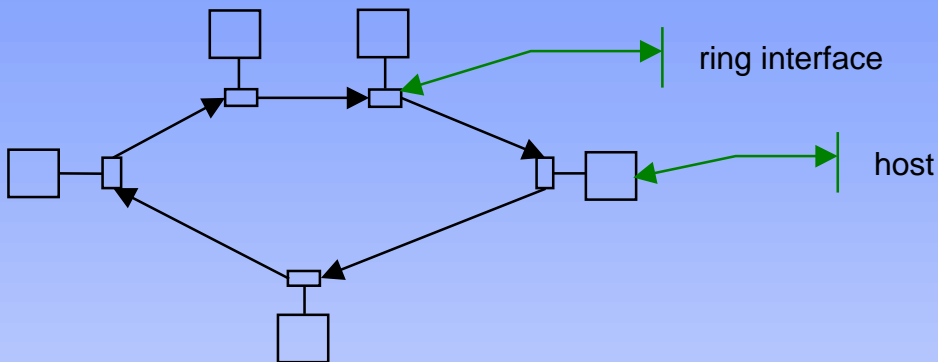
- Tree:



- data moves in both directions on a branch, “headend” retransmits on other branches
- each branch is a bus (i.e. a bus is just a single-branch tree)
- size of branches (number of headends) is a major cost factor
- analog broadband (cableco) for MAN trees
 - problem: analog signalling is unidirectional, required FDM for bidirectional

Ring

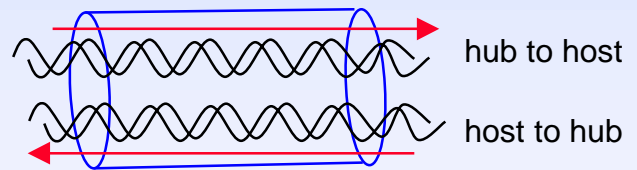
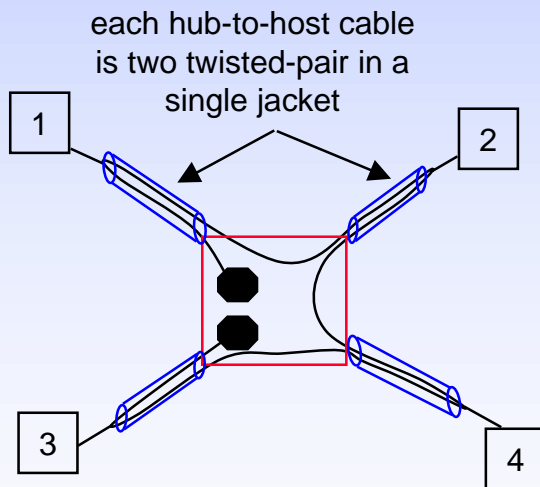
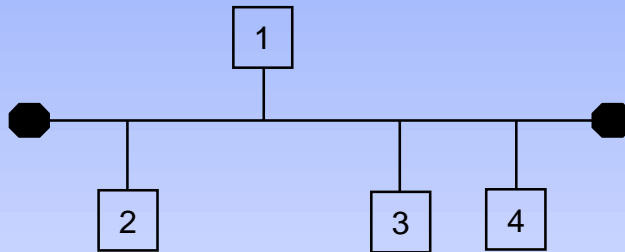
- Ring is formed from a sequence of point-to-point connections



- Connections are really between “ring interface” devices (Tanenbaum)
- Ring interface reads each bit and retransmits it (listen mode)
- Obviously each link is not “shared” in the bus broadcasting sense
 - “shared” because all links must be used by all stations for arbitrary point-to-point
- Data flow is unidirectional (links all go in the same direction)

Star: layout or topology?

- A star-shaped network (all hosts joined at a central location) can be considered just a wiring layout or a full-fledged topology
- A bus can be transformed into a star:



- A ring can be similarly transformed

...star: layout or topology, 2

- In the preceding simple case, the device at the centre is passive (doesn't contain any active circuitry or software)
 - called a *hub*, also, wiring hub, repeater
 - hubs operate at the physical layer
 - doesn't really change the topology, just a layout convenience
- Star/hub wiring can help with network management
 - failure in linear bus wiring hard to find, affects every host
 - failure in star wiring affects just the one host
- But, the hub becomes a single point of failure
 - if it fails, no communication anywhere
- A *switch* is a device that is positioned like a hub, but is entirely different
 - a switch operates at the data link layer;
discussion to follow

Commercial MAC implementations

- MAC sublayer is responsible for managing contention and resolving collisions
 - bus: simultaneous transmissions travel along the medium and collide
 - ring: two stations injecting data into the ring will collide
- There is an almost universal family of MAC standards and protocols
 - the 802.x series standardized by IEEE
 - 802.2: LLC – logical link control; equivalent to upper half of OSI DL
 - 802.3 and above: lower half OSI DL and physical layer
 - common addressing scheme (6 byte) and other commonality where feasible
 - called *MAC addresses*
 - addresses are assigned to hardware by manufacturer (6 byte = 48 bit $\Rightarrow 2^{48} \approx 3 \times 10^{14}$ (30 trillion))

Ethernet 802.3

- Ethernet™ is a bus topology that uses Carrier Sense Multiple Access with Collision Detection (CSMA/CD):
 - standardized as IEEE 802.3
 - basic operation
 - host waits for idle bus (no carrier)
 - transmits; listens for collision
 - if collision, stop
 - wait random amount of time, repeat
- Cabling options (historic interest only):
 - 10Base5: “thick Ethernet”, 10Mbps, baseband coaxial, 500 m/segment
 - 10Base2: “thin Ethernet”, 10Mbps, baseband coaxial, 200 m/segment
 - 10Base-F: 10Mbps, fiber, 2km/segment
 - 1Base5: 1Mbps, baseband twisted-pair, 250 m/segment
 - 10Broad36: 10Mbps, broadband coaxial, 3.6km/segment

... Ethernet, 2

- Standard use cabling:
 - 10Base-T: 10Mbps, baseband twisted-pair, 100 m/segment; hub with star layout
 - use is waning, legacy systems
 - some 10Base2 is still used (older legacy)
- Most common new installations are “fast” Ethernet, 802.3u
 - 100 Mbps
 - twisted-pair with central hub or switch
 - cabling options:
 - 100Base-T4: uses 4 twisted pairs; 100 m/segment
 - 100Base-TX: uses 2 “Category 5” (high quality) twisted pair; 100 m/segment
 - 100Base-TF: uses 2 fiber-optic strands; 2000 m/segments
 - note that distances are shorter with copper, inter-building can be a problem

... Ethernet, 3

- In Ethernet, choice of wait period after collision is “random”, but adjusts to number of stations
 - Ethernet is classified as a *random-access dynamic channel allocation* protocol
 - time period increases as number of consecutive collisions increase
 - called the *backoff* period; *exponential backoff* doubles the wait time after each
- 802.3 works well for lightly-loaded networks
 - as traffic increases, probability of collision increases
 - traffic usually increases as the number of hosts connected to the medium increases
- Ethernet has the operational characteristic that it scales poorly

... Ethernet, 4

- Solutions for overloaded 802.3:
 - more speed
 - undesirable price-performance
 - fewer hosts
 - partition the hosts into smaller groups, join groups together in a *switch*
 - the medium is partitioned
 - each partition has fewer hosts
 - fewer hosts means fewer collisions (probably);
 - the *collision domain* has been reduced in size and the probability of collision is reduced
- An 802.3 switch is a device at the centre of a star (same position as a hub) that performs *collision domain partitioning* and improves Ethernet performance

... Ethernet, 5

- Switch operation is a significant product differentiator:
 - there are many product configurations, options and operating characteristics
- Typically:
 - store-and-forward frames at the data-link layer
 - variation: “cut-through switching” starts sending a frame to its destination before it’s fully received
 - may support full duplex or only half duplex
 - can handle combinations of 10 and 100 Mbits/sec hosts (and even 1Gbits/sec)
 - may have from 4 to 96 host interfaces of various speeds
 - are implemented with special-purpose hardware and software
 - may be as simple as a collection of hubs interconnected with router-like software, or exceedingly complex

... Ethernet, 6

- Switches can be arranged in a hierarchy
 - two-layer hierarchy of passive hubs connected to small switches can be quite cost effective
 - there are many product configurations, options and operating characteristics
- Ethernet *bridges* perform a similar task to switches
 - a bridge in the generic sense is a device that connects two networks together at the datalink layer
 - this is more or less what a switch does, so the distinction is quite unclear
 - WAN bridges are generally devices at each end of a point-to-point link (two ports)
 - Kurose & Ross simply suggest that Ethernet bridges have fewer ports

... Ethernet, 7

- Both switches and bridges are designed to be transparent (plug & play)
 - they have to learn about the topology and connected devices by inspecting frames and building tables or databases
- Extra confusion: joining LANs is also something that a router can do
 - consider IP subnetting: dividing an IP network into smaller physical networks
 - often done for network administrative autonomy, but could be done for load management
 - each network has a router to the outside world
 - routers are definitely **not** plug & play

Other Ethernet & LAN standards

- Really fast Ethernet:
 - “gigabit Ethernet” (1,000 Mbits/sec) being deployed
 - 802.3z fiber
 - 802.3ab 1000Base-T; Cat5 twisted (4 pairs)
 - not common on the desktop – yet
 - on the drawing board: 802.3ae 10Gbits/sec
- Existing fiber technology is capable of handling transmission rates of 50,000 Gbits/sec (50 Terabits/sec) easily
 - problem is electro-optical conversion
 - research into pure optical computers is crucial

... other standards, 2

- Wireless LANs: 802.11
 - designed to replace wires in LANs
 - traditional deployment has hosts with transponders organized around *access points* (base stations)
 - access points can be connected together with traditional technologies to form larger networks & “wandering” support
 - 802.11 also permits hosts with transponders to form *ad hoc* groups without a defined access point
 - e.g. laptops in a meeting room
 - 802.11 original definition: 1 Mbits/sec and 2 Mbits/sec
 - speeds based on availability of radio frequency spectrum, transmitter power rules, ITU-R bureaucracy
 - faster standards are now available (up to 11 Mbits/sec)
 - an infrared version is also defined

... other standards, 3

- Wireless Personal Area Networks: 802.15
 - WPANs
 - equivalent proprietary systems:
 - Bluetooth: Ericsson *et al*
 - HomeRF: developed Shared Wireless Access Protocol (SWAP), 1998;
 - interoperability of 802.11 devices with the ever-growing set of PDAs, pagers, WAP devices, PCS phones, etc.
 - differing goals from 802.11:
 - cheaper
 - low power consumption
 - simpler devices, lower data rate requirements (initially at least)
 - initial data rate requirements up to 1Mbit/sec

Other 802.x LAN technologies

- Token bus; 802.4
- Developed by GM and others for factory automation applications
- Goals: response-time predictability and cabling robustness
- Physical: broadband coaxial bus
 - 1, 5, 10 Mbps
 - moderate bandwidth, long distances
- MAC contention is handled by avoidance with turn-taking protocols
 - *token passing*: a sender must possess a token before sending
 - no token, no sending!
 - may hold the token for predetermined time period, then must let another host have it
 - token is just circulated between hosts if no host wants to send

... other 802.x, 2

- Token ring; 802.5
- Similar goals and ideas as token bus
 - predictable response times, even in heavy traffic
- Physical: sequence of point-to-point twisted-pair connections, forming a physical ring
 - 4, 16, 100 Mbits/sec
 - wired in a star with a hub
- MAC contention handled by token passing (same ideas as token bus)
- Token ring was popularized by IBM early in the development of the PC
 - gained significant following in IBM corporate and enterprise clientele
 - not for general use now (still applicable for guaranteed-response situations)

LLC – logical link control

- Standardized as 802.2; based on HDLC
- “Semi-layer” above other 802 standards
- Provides typical DL functions like flow-control, error checking
- Three defined services
 - connection-oriented
 - unacknowledged connectionless
 - acknowledged connectionless
- LLC not used in IP networks (IP uses MAC sublayer directly)

Other LAN technologies

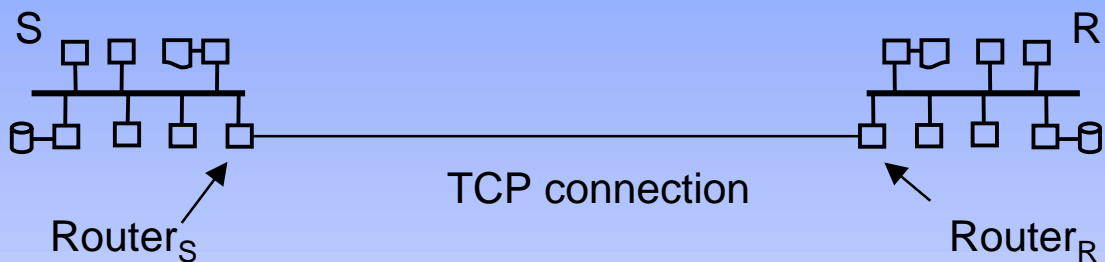
- FDDI: Fiber Distributed Data Interface
 - ring topology with token-ring MAC
 - fiber-optic media; 100Mbps
 - uses two rings, one “clockwise”, other “counterclockwise”
 - redundancy for fault tolerance
 - can form one ring out of two semi-rings, if necessary
 - each station has hardware to join rings
 - typical uses: backbone LAN (connecting smaller LANs)
 - up to 200 km with 500-1000 stations

... other LAN technologies, 2

- ATM
 - early in its history, ATM was considered a potential replacement for high-speed Ethernet to the desktop
 - cheap highspeed Ethernet and TCP/IP ended the debate
 - ATM had a APIs defining how application would interface directly with ATM's uppermost layer (roughly a network layer)
 - LANE: emulate 802.3 & 5 over ATM
 - didn't happen, and probably won't
 - ATM is a very strong presence in WAN backbone technologies, especially with the developments of IP directly on ATM
- But, there is also work on running packet-switching directly on SONET (confusingly called POS: Packet Over SONET)
 - recall that SONET is synchronous, so statistical packet multiplexing was not the intended use of SONET

TCP/IP and 802 LANs

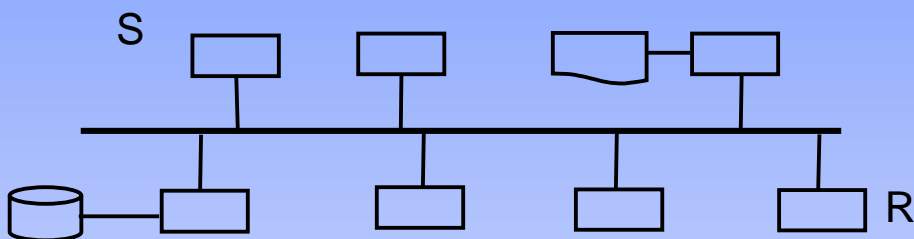
- TCP/IP is an internetworking system
 - common use is to connect 802.x LANs
- Conceptually easy:



- Problem: host S sends to R, using TCP/IP address (IP address and port)
 - assume that S knows how to get to its router (Router_S)
 - how does the receiving router Router_R know which host is R?
 - within Router_R's LAN, R is identified by its MAC address
 - need translations between MAC addresses and IP addresses

... TCP/IP and 802 LANs, 2

- An equivalent problem exists for a single LAN that uses TCP/IP for its transport and higher layers



- Within a single LAN, no internetworking is required: hosts only need the LAN MAC addresses
 - if S and R use “internet” names and addresses to identify each other, translation to MAC address is required
 - DNS is not sufficient, since it translates names to IP addresses, not MAC addresses

... TCP/IP and 802 LANs, 3

- LANs need not use TCP/IP addressing internally
 - PCs, Apple both have well-established LAN protocol suites that are not TCP/IP
 - Windows NT networking: NT domains, workgroups, NetBEUI
 - LocalTalk/AppleTalk
 - Novell
 - Banyan
 - all of these systems handled the naming and addressing of hosts with proprietary protocols
- In many LANs, there are two independent sets of upper-layer protocols that “meet” at the MAC layer
 - or, more commonly now, use IP as a datagram service for proprietary transport & higher

ARP – address resolution protocol

- So, how to solve the IP vs LAN addressing problem?



- Example:
 - S sends to R with IP address (obtained from DNS usually)
 - if S knows only the IP address of Router_S, use ARP to get its MAC address
 - Router_S uses normal internet routing to get to Router_R (e.g. by extracting the network address from Router_R's IP)
 - Router_R receives data for a host on its network, but only knows the IP address
 - Router_R uses ARP to discover the MAC address corresponding to R's IP address
- ARP is a TCP/IP host-to-network (link) layer protocol

... ARP, 2

- ARP operation (for Router_R):
 - Router_R broadcasts an “ARP request” on its LAN:
 - “would the owner of IP address a.b.c.d please respond”
 - LAN broadcasting is easy
 - station R responds, replies with an “ARP reply”
 - “I am a.b.c.d; my MAC address is z-y-x-w-v-u”
 - Router_R then sends the data directly to R using its MAC address
- For efficiency, ARP information can be cached by routers (and hosts)
- Proxy ARP: a host that answers ARP requests for many IP addresses
 - lets the host act as an IP “funnel”
 - e.g.: dialup terminal server answers ARP request for each connectee

Reversed ARP

- The opposite problem: given a MAC address, how to determine the IP address?
 - in a workstation with secondary storage (disk), use a configuration file or similar repository
 - what about stations that have no such storage (e.g. X terminals, “network computers”)?
- Solution: reversed ARP – each LAN has a server that can respond to broadcasts:
 - “ my MAC address is z-y-x-w-v-u. does anyone know my IP address?”
 - RARP server has a database of MAC–IP mappings: it responds directly with:
 - “your IP address is a.b.c.d”

... reversed ARP, 2

- Problem:
 - LAN broadcasts cannot be routed across networks
 - need a RARP server on each network
- Bigger problem:
 - in general, stations need need more information than an IP address
 - other configuration information
 - boot image
- Solution: BootP – Boot Protocol
 - BootP is an extension of the RARP idea using routable UDP datagrams instead of LAN broadcasting
- BootP can be used to get other information necessary to boot the station
- DHCP (Dynamic Host Configuration Protocol) is an extension of BootP that allows reusable IP addresses

DHCP

- DHCP is designed to make IP address and other configuration information (e.g. host names, subnet masks, DNS addresses) “plug & play”
- A DHCP server can be set up to assign information dynamically, or can reserve based on MAC address
 - DHCP-provided information has an associated *lease time*
 - when a lease expires, the host must reacquire configuration information
 - could be the same information, or different
- DHCP is most applicable to general workstations and especially mobile hosts
 - servers need to be at known addresses
 - DHCP can be used with servers; generally uses reservations of known addresses (management convenience)

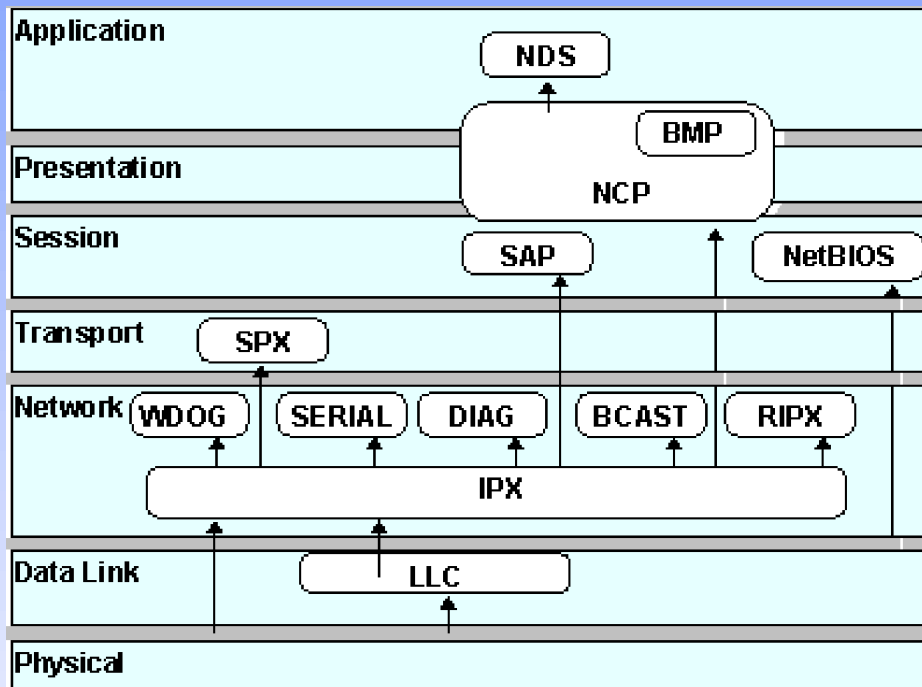
... DHCP, 2

- Hosts that use DHCP are not configured with a DHCP server address
 - when host starts, it broadcasts a *DHCP discover* message to see if any DHCP servers are available
 - host choose among responses to obtain information
- DHCP is now common for ISPs to assign configuration information customers
 - even though DSL and cableco offerings are “always on”, DHCP is still used for management purposes
- Most “Internet sharing” (NAT) devices also have a DHCP server to provide information to the interior (private IP)
- DHCP does not coordinate with DNS to create DNS entries assigned addresses
 - newer version of DHCP & DNS do

A non-TCP/IP LAN

- Novell Netware: a LAN based on 802.x standards, without TCP/IP

- from <http://www.protocols.com/pbook/novel.htm>



- 4 or 5 layers:
 - layers don't align exactly with OSI, but close

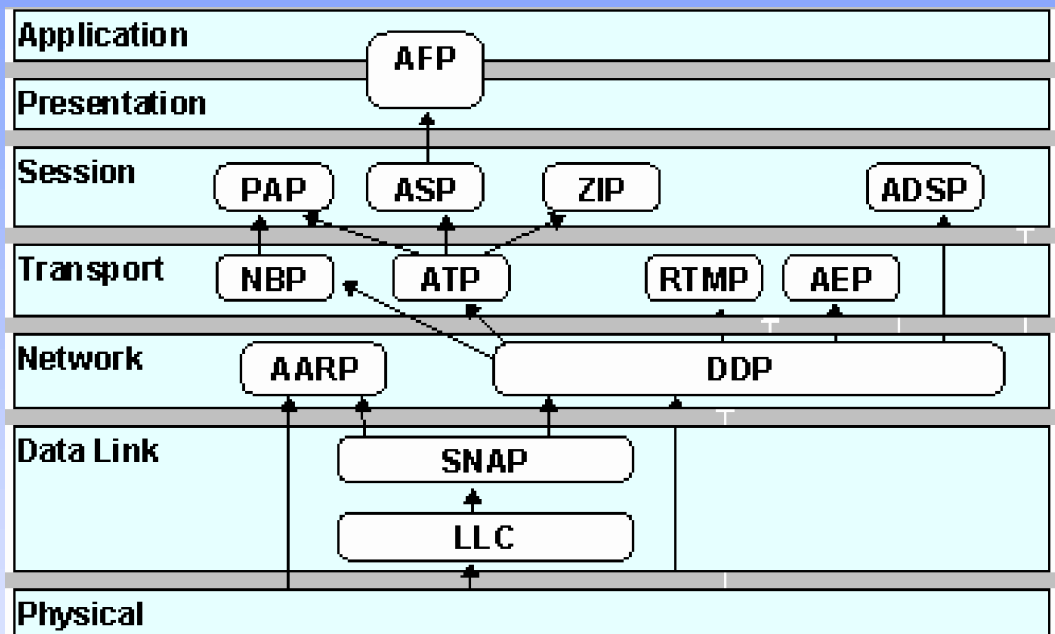
... Novell netware, 2

- Application to session layers:
 - NCP: NetWare Core Protocol; connection-oriented++
 - NDS: NetWare Directory Services
 - SAP: Service Advertising Protocol
- Transport:
 - SPX: sequenced packet exchange; plain connection-oriented
- Network
 - IPX: internet packet exchange; unreliable connectionless, 12-byte addresses
 - similar in philosophy to IP: internetworking, routable
- DL and physical
 - from 802.x

Another LAN – AppleTalk

- AppleTalk from Apple Corp. protocol suite

– from <http://www.protocols.com/pbook/appletalk.htm>



- Physical layer can be LocalTalk, or 802.3
- Internetworking, routable
- **AARP** - AppleTalk Address Resolution Protocol; **DDP** - Datagram Delivery Protocol; **RTMP** - Routing Table Maintenance Protocol; **AEP** - AppleTalk Echo Protocol; **ATP** - AppleTalk Transaction Protocol; **NBP** - Name-Binding Protocol; **ZIP** - Zone Information Protocol; **ASP** - AppleTalk Session Protocol; **PAP** - Printer Access Protocol; **ADSP** - AppleTalk Data Stream Protocol; **AFP** - AppleTalk Filing Protocol

Yet another LAN – Windows NT

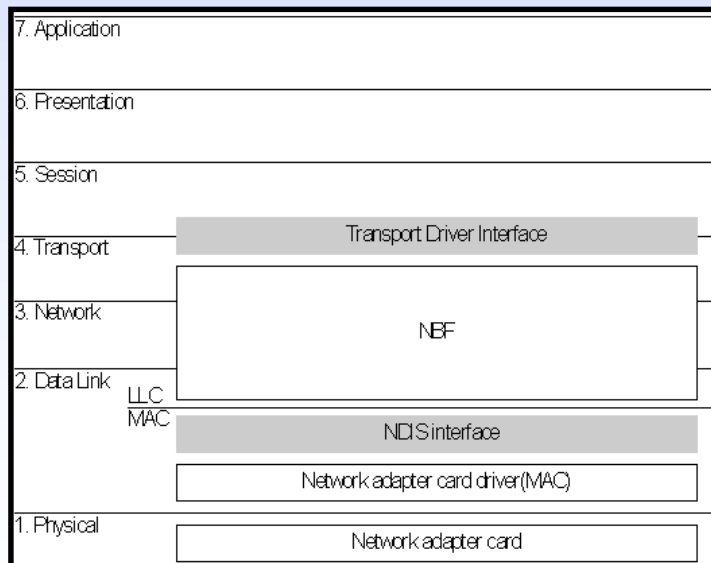
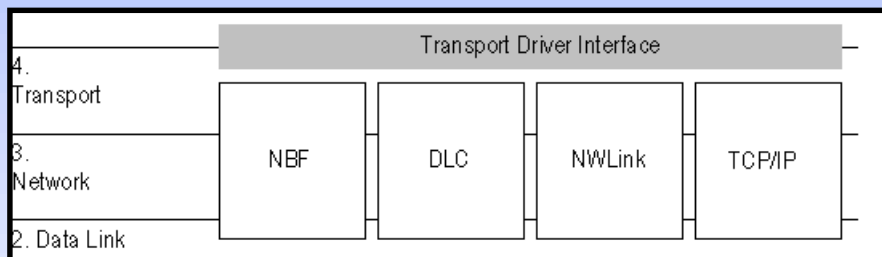
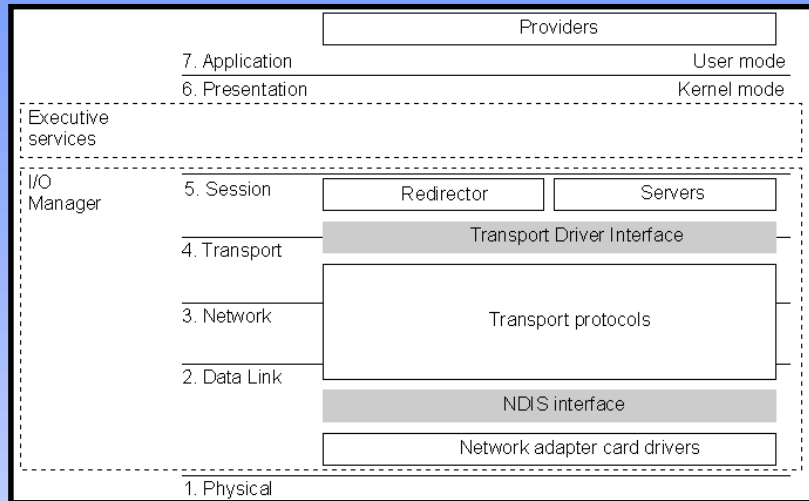
- Windows NT networking is a session-to-application suite of protocols inherent in Window NT and related operating systems
 - supports client-server and peer application models
 - file browsing and sharing, printing, authentication
 - provides TCP/IP socket services
 - provides versions of TCP/IP applications an protocols
 - DNS, DHCP, SMTP, FTP
 - not the same behavior as the Unix originals, which is a source of irritation
 - some interoperability with AppleTalk & Novell application layers

... yet another LAN – Windows NT, 2

- Can use any of four transports
 - Novell IPX/SPX
 - TCP/IP
 - including all of the standard TCP/IP protocols: ARP, PPP, DHCP clients
 - NBF (NetBEUI Frame)
 - NetBEUI/NBFI is primarily a network layer process, with some reach into transport
 - CO and CL services, sliding window for flow and congestion control
 - uses MAC address directly, **not routable**
 - provides services to the NetBIOS (Network Basic I/O Services) layer; which is a transport/session layer
 - DLC – Data Link Control (from the OSI HDLC, SNA SDLC, X.25 LAP & LAPB family of datalink protocols)
 - used for mainframe, network printers
 - contains an 802.2 LLC interface for driving network interface cards

... yet another LAN – Windows NT, 3

- Pictures:



Summary

- LANs are built from collections of hosts connected to shared media
 - requires MAC – medium access control
 - broadcasting to all hosts is easy
- Various hardware implementations exist, but IEEE 802.3 Ethernet CSMA/CD is ubiquitous and the *de facto* standard
 - it has problems under loaded conditions, but ever-increasing data rates and switching technology helps to extend the life of the technology
- TCP/IP was never intended to be the lower-layer protocols within a LAN, so MAC-to-IP address translation is provided by ARP
- Proprietary non-TCP/IP LANs exist and are generally all built on 802.x